

LECTURE 2 || 2 : SYN. MONOIDAL CATEGORIES and 3 FUNCTOR FORMALISM

20. MOTIVATION

Recall: WANT a "6-FUNCTOR FORMALISM" for \mathcal{C} category w - x -
 $E \subseteq \text{Hoph}(\mathcal{C})$ stable under $\dashv\dashv$
"EXCEPTIONAL MORPHISM" pullback
isomorphs

Main Example: • $\mathcal{C} = \text{loc. compact Hausdorff top spaces}$

$E =$ "compactifiable maps"
i.e. maps which write as

$$f = \bar{f} \circ j \quad \begin{array}{l} \swarrow \text{open immersion} \\ \uparrow \\ \text{proper maps} \end{array}$$

- $\mathcal{C} = \text{QCohent schemes}$
 $E = \text{separated morphisms of finite type}$

where where A 6 FUNCTOR FORTAUST should be.

- EASY to FORTAUST
 (it is a functor
 $C^{op} \rightarrow \text{cat of s.m. cat}$)
- 1) $C \ni X \mapsto D(X) \in \text{Cat}$
 - 2) $(D(X), \otimes)$ s.m. structure
 - 3) $\forall f: X \rightarrow Y$, a FULLBACK functor $f^*: D(Y) \rightarrow D(X)$
 compatible w/ $-\otimes-$, $()^*$ functorial

- 4) $\forall f: A \rightarrow B$ in E , an EXCEPTIONAL pushforward $f_!: D(X) \rightarrow D(Y)$
 $()_!$ functorial on E , BASE CHANGE, PROJECTION FORMULA

This is the HAND PART!

- 5) RIGHT ADJOINTS: $-\otimes- \rightarrow \text{Hom}(-, -)$
 $f^* \rightarrow f_*$
 $f_! \rightarrow f^!$

\Rightarrow SOLUTION : Define a sym. mon. cat $\text{Corr}(C, E)^{\otimes}$
 (to have $\otimes, f^*, f_!$)
 so that (i) \rightarrow (ii) \iff LAX S.M. FUNCTOR
 $\text{Corr}(C, E)^{\otimes} \rightarrow \text{Cat}$

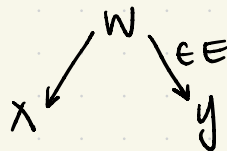
31. The Category $\text{Corr}(C, E)$

DEF (last time) C locat w/ finite limits
 $E \subseteq \text{Morph}(C)$ stable under $- \circ -$, pullback, \geq isos

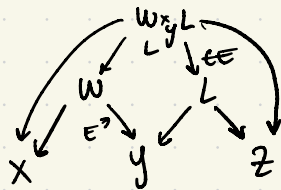
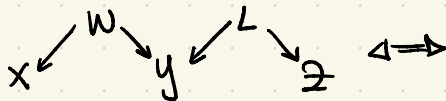
$\text{Corr}(C, E)$ should be the category with:

- $\text{Ob}(\text{Corr}(C, E)) = \text{ob}(C)$

- morphisms: $X \xrightarrow{\phi} Y \iff$



- composition:



Also would like a s.m. structure

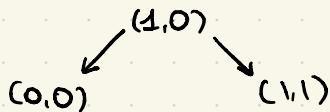
where $X \otimes Y = X \times Y$

Now: $\text{Corr}(C, E)$ as an locat.

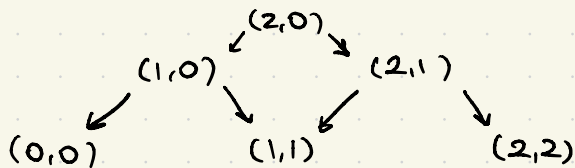
What are s.m. categories and $\text{Corr}(C, E)^{\otimes}$ s.m. struct

DEF • $\forall n$, $(\Delta^n)_+^2 \subseteq (\Delta^n)^{\text{op}} \times \Delta^n$ subsimpl set spanned by (i,j) w/ $i \geq j$

n=1



n=2

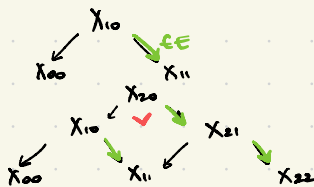


ms $(\Delta^\bullet)_+^2 : \Delta \mapsto \text{sSets}$ (COSIMPICAL DCATEGORY)

• $\text{Corr}(CE)_\bullet \subseteq \text{Fun}^\infty((\Delta^\bullet)_+^2, \mathbb{C})$ subsimplical set where

$f \in \text{Corr}(CE)_n \iff \begin{cases} \text{ALL ARROWS DOWN RIGHT} \in E \\ \text{ALL SMALL SQUARES ARE CANTESIAN} \end{cases}$

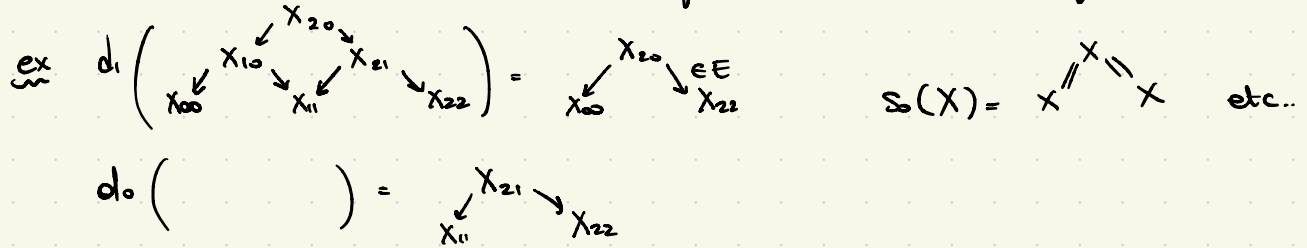
n=1



n=2

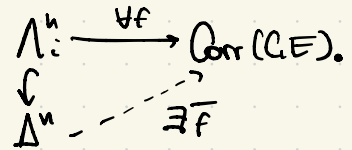
$x_{ij} \in \text{ob}(\mathbb{C})$

Rmk Well defined bc $f \in \text{Corr}(C, E)_n \Rightarrow$ dif $\in \text{Corr}(C, E)_{n-1}$ $\forall i$
 $s_j f \in \text{Corr}(C, E)_{n+1}$ $\forall j$



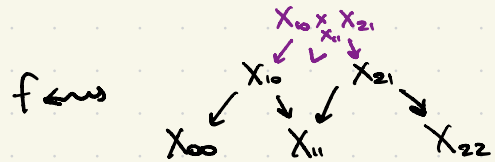
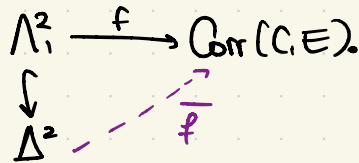
PROP $\text{Corr}(C, E)$ is an ∞ -category, i.e. HORN FILLING CONDITION

i.e. $\forall n \geq 2, \forall 0 \leq i \leq n$



PROOF [lemma 6.1.2 of [L212a]]

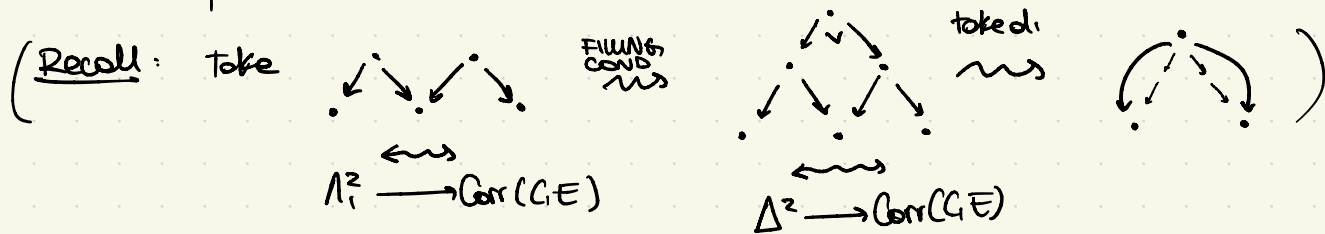
For $n=2, i=1$



\bar{f} can be obtained via



Rmk In particular, we know that composition in the most $\text{Corr}(C, E)$ corresponds to what we wanted



2.2. SYNTH. MONOIDAL ∞ -CATEGORIES

Goal | way to be able to define $\bigotimes_{i \in I} X_i$ $\forall I$ finite set, $X_i \in \text{ob}(C)$ |

CLASSICAL DEF of s.m. 1-cat (w/ \otimes, \dashv : $C \times C \rightarrow C$)
 $1 : * \rightarrow C$
 isos: $x \otimes (y \otimes z) \simeq (x \otimes y) \otimes z$
 $x \otimes y \simeq y \otimes x$
 $1 \otimes x \simeq x$
 + coherence (pentagon + hexagon axioms)

\rightsquigarrow HARD TO GENERALIZE
 (HIGHER COHERENCES!)

\rightsquigarrow Reformulation: "A S.M. 1-cat = COMMUTATIVE MONOID in (Cat, \times) "

What is a COMMUTATIVE MONOID in $\mathbf{Cat}_{\text{fin}}$? ← assoc of monoids

DEF Fin. category w/ objects FINITE POINTED SETS, denoted by $\langle I \rangle = I \cup \{*\}$, I finite set
 morphisms \rightsquigarrow mops of finite pointed sets

Special objs $\langle n \rangle = \{1, \dots, n\} \cup \{*\}$

Some special mops • $\forall I$ fin set, $\forall i \in I$

$$\langle I \rangle \xrightarrow{e_i} \langle i \rangle$$

$$j \longmapsto i \quad \text{if } j=i$$

$$* \quad \text{otherwise}$$

"projections"

$$\bullet \beta: \langle 2 \rangle \longrightarrow \langle 1 \rangle$$

$$* \longrightarrow *$$

$$1 \longrightarrow 1$$

$$2 \longrightarrow 1$$

"multiplication"

DEF | A SYM. MONOCATEGORY is a functor $X: \mathbf{Fin}_{\text{fin}} \longrightarrow \mathbf{Cat}_{\text{fin}}$ s.t. $\forall I$ fin set

the map $X(\langle I \rangle) \xrightarrow{\sim} \prod_{i \in I} X(\langle i \rangle)$ induced by $\langle I \rangle \xrightarrow{e_i} \langle i \rangle$

is an equivalence.

("Segal condition")

Good def bc

• objectwise, X determined by $X(\langle 1 \rangle) \in \text{Cat}$ $D = X(\langle 1 \rangle)$

• TENSOR PRODUCT: should be $-\otimes- : D \times D \rightarrow D$

↳ it is the image of $\beta : \langle 2 \rangle \rightarrow \langle 1 \rangle$!!

Indeed $X(\beta) : X(\langle 2 \rangle) \xrightarrow{\cong} X(\langle 1 \rangle) \times X(\langle 1 \rangle) = D \times D \rightarrow X(\langle 1 \rangle) = D$

and it is: • SYMMETRIC because

$$\begin{array}{ccc} \langle 2 \rangle & \xrightarrow{(12)} & \langle 2 \rangle \\ \beta \searrow & & \swarrow \beta \\ & \langle 1 \rangle & \end{array} \Rightarrow \begin{array}{ccc} D \times D & \xrightarrow{-\otimes-} & D \\ \downarrow (12) & & \nearrow -\otimes- \\ D \times D & & \end{array}$$

• ASSOCIATIVE because

meaning $a \otimes b \simeq b \otimes a$ (All higher coherences encoded)

$$\begin{array}{ccc} \langle 3 \rangle & \longrightarrow & \langle 2 \rangle & \xrightarrow{\beta} & \langle 1 \rangle \\ \begin{array}{c} 1 \\ 2 \\ 3 \end{array} & \longrightarrow & \begin{array}{c} 1 \\ 2 \end{array} & \longrightarrow & 1 \end{array}$$

in Fin.

$$\begin{array}{ccc} \langle 3 \rangle & \longrightarrow & \langle 2 \rangle & \xrightarrow{\beta} & \langle 1 \rangle \\ \begin{array}{c} 1 \\ 2 \\ 3 \end{array} & \longrightarrow & \begin{array}{c} 1 \\ 2 \end{array} & \longrightarrow & 1 \end{array}$$

$$\Rightarrow \begin{array}{ccc} D^{\times 3} & \xrightarrow{(id, -\otimes-)} & D^{\times 2} \\ \downarrow (\otimes, id) & & \downarrow \otimes \\ D^2 & \xrightarrow{-\otimes-} & D \end{array}$$

while it is MORE COMPLICATED if C has another sm strict ...

• UNITAL because $X(\langle \emptyset \rangle) \simeq * \rightarrow X(\langle 1 \rangle) = D$

Rmk If C is a cat w/ finite pools, a commutative monoid in C is $\text{Fin} \rightarrow C$ w/ Segal condition

PRACTICAL PROBLEM: Hard to write functors into Cat

⊆
Already at 1-categorical level, where usually one has PSEUDOFUNCTORS i.e. $X(f \circ g) \cong X(f) \circ X(g)$
only iso
no longer a pbm in Cat but these higher coherences

Pbm in Cat solved by GROTHENDIECK

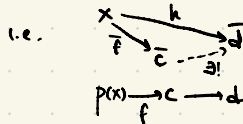
Thm | There is an EQUIVALENCE of 2-CATEGORIES

$$\text{PseudoFunctors}(C, \text{Cat}) \cong \text{coFib}(C)$$

COANCESTRAL FIBRATIONS over C

where $p: E \rightarrow C$ is a coCartesian fibration if $\forall p(x) \xrightarrow{f} c$ morph in C

\exists lift $x \xrightarrow{\bar{f}} \bar{c}$ ($p(\bar{f}) = f$)
 which is INITIAL amongst lifts } "COANCESTRAL LIFT"



• The correspondence goes as:

$$\begin{array}{ccc} \text{PstFun}(C, \text{Cat}) & \longleftarrow & \text{coFib}(C) \\ F: C \rightarrow \text{Cat} & \longleftarrow & (E \xrightarrow{p} C) \\ c \rightarrow p'(c) =: E_c & & \end{array}$$

$$\begin{array}{ccc} \text{PstFun}(C, \text{Cat}) & \longrightarrow & \text{coFib}(C) \\ (F: C \rightarrow \text{Cat}) & \longmapsto & \left(\int_F C \xrightarrow{p} C \right) \end{array}$$

where $\int_F C$ is a category where: • Obj's = $(c \in C, x \in \text{Ob}(F(c)))$

• Morphisms $(c, x) \rightarrow (d, y)$ is $\begin{cases} f: c \rightarrow d \text{ in } C \\ \begin{matrix} f_c \xrightarrow{Ff} f_d \\ \downarrow \psi \\ x \rightarrow Ff(x) \end{matrix} \end{cases}$ and $\phi: Ff(x) \rightarrow y$

$\int_F C \xrightarrow{F} C$ forgetful functor.

There is an ANALOGOUS RESULT for ∞ CATEGORIES

DEF $p: E \rightarrow C$ b/ ∞ cats is a CoCartesian fibration if

p inner fibration

$\forall p(x) \xrightarrow{f} c \quad \exists x \xrightarrow{\bar{f}} \bar{c}$ lift which is INITIAL
(ie. space of --- is contractible)

"locally coCartesian lift"

locally coCart. lifts stable under composition

Rmk $p: E \rightarrow C$ coCart fib of 1-cat

$\Leftrightarrow N(p): N(E) \rightarrow N(C)$ coCart fib of ∞ cat

Thm [Straightening-
Unstraightening,
Lurie] There is an EQUIVALENCE of $(\infty, 1)$ -categories

$$\mathrm{Fun}^{\infty}(C, \mathrm{Cat}_{\infty}) \simeq \mathrm{coCartFib}(C)$$

where on objects it works as the 1-categ case

Rmk Ift would be nat equiv of $(\infty, 2)$ -cat. But we take $\mathrm{Fun}^{\infty}(C, \mathrm{Cat}_{\infty})$ ONLY INVERTIBLE NATURAL TRANSFORM

\Rightarrow in the ∞ cat of $\mathrm{coCartFib}$ we restrict to
$$\begin{array}{ccc} E & \xrightarrow{F} & E' \\ P \downarrow & & \downarrow P' \\ C & & C \end{array}$$
 st F preserves loc. coCart . lifts

DEF" A S.T. COCAT is a COCARTESIAN FIBRATION $C^{\otimes} \xrightarrow{P} \mathrm{Fin}$. s.t. $\forall I \in \mathrm{Fin}$ set,

the nat map $(e^{\otimes})_{\langle I \rangle} \xrightarrow{\sim} \prod_{i \in I} (e^{\otimes})_{\langle i \rangle}$ induced by coCart lifts of $\langle I \rangle \xrightarrow{e_i} \langle i \rangle$

is an equiv of ∞ categories

We call $(e^{\otimes})_{\langle i \rangle} =: \mathcal{C}$ the UNDERLYING ∞ CATEGORY of P .

Hyponotic $X: \text{Fin.} \rightarrow \text{Cat}_{\text{co}}$ s.m. 1-cat, $C = X(\langle 1 \rangle)$

\leftarrow straight-unst $p: C^{\otimes} \rightarrow \text{Fin.}$

$$\text{Ob}(C^{\otimes}) = \{ (c_i)_{i \in I} \mid c_i \in C, I \in \text{fin sets} \}$$

$$C^{\otimes}((c_i)_{i \in I}, (d_j)_{j \in J}) = \left\{ \alpha: \langle I \rangle \rightarrow \langle J \rangle \text{ in Fin. and } \begin{cases} \forall j \in J & f_j: \bigotimes_{i \in \alpha^{-1}(j)} c_i \rightarrow d_j \end{cases} \right.$$

$-\otimes-$ is the cotensorial left of $\beta: \langle 2 \rangle \rightarrow \langle 1 \rangle$

Prop Def" of s.m. coact is apparently not self dual, but via def' and $\text{Cat}_{\text{co}} \simeq \text{Cat}_{\text{co}}$ it is.
 $C \rightarrow C^{\text{op}}$

SPECIAL CASE

C w/ finite prods $\Rightarrow C^{\text{op}}$ w/ finite coprods

[2.6.1.4 HA
2.6.1.8]

and $\exists!$ $(C^{\text{op}})^{\cup} \rightarrow \text{Fin.}$ sym. monoidal category where $-\otimes-$ = $-\cup-$

C w/ finite prods $\Rightarrow \exists!$ $C^{\times} \rightarrow \text{Fin.}$ s.m. structure on C where $-\otimes-$ = $-\times-$

Hyponotic $C = \text{Cat}_{\text{co}}$

FINAL NOTION NEEDED : LAX S.M. FUNCTOR

1-Categorical: $\phi: (C, \otimes) \rightarrow (D, \otimes)$ w/ natural maps $\bigotimes_{i \in I} \phi(x_i) \rightarrow \phi(\bigotimes_{i \in I} x_i) \quad \forall I$

\iff want existence of natural transformations

$$\begin{array}{ccc} C^{\times I} & \xrightarrow{\phi^{\times I}} & D^{\times I} \\ \otimes \downarrow & \swarrow & \downarrow \otimes \\ C & \xrightarrow{\phi} & D \end{array}$$

$$\begin{array}{ccc} F_C: \text{Fin.} \rightarrow \text{Cat} & & F_D: \text{Fin.} \rightarrow \text{Cat} \\ F_C(I) \xrightarrow{\phi_I} F_D(I) & & \\ \downarrow \pi & \xrightarrow{\pi \phi} & \downarrow \pi \\ \prod F_C(\langle i \rangle) & & \prod F_D(\langle i \rangle) \\ \otimes \downarrow & & \downarrow \otimes \\ F_C(\langle i \rangle) \xrightarrow{\phi_{\langle i \rangle}} F_D(\langle i \rangle) & & \end{array}$$

- i.e. $\left\{ \begin{array}{l} \bullet \text{ pseudonatural transformation (not necessarily invertible)} \\ \bullet \phi_I \simeq \prod_{i \in I} \phi_i \end{array} \right.$

DEF
(LAX S.M. FUNCTOR)
b/ ∞ CATS

STRAIGHT-UNSTRAIGHT

$$\begin{array}{ccc} C^{\otimes} & \xrightarrow{\tilde{\phi}} & D^{\otimes} \\ \downarrow p & & \downarrow p' \\ \text{Fin} & & \end{array} \quad \tilde{\phi}: C^{\otimes} \rightarrow D^{\otimes} \text{ over Fin}$$

st $\tilde{\phi}$ preserves loc cocart lifts of $\ell_i: \langle I \rangle \rightarrow \langle i \rangle$

(then ϕ s.m \iff not transf are invertible $\iff \tilde{\phi}$ preserve all (loc cocart) lifts)

43. S.M. STRUCTURE on $\text{Corr}(C, E)$ for (C, E) geometric context

Goal: Define $\text{Corr}(C, E)^{\otimes} \rightarrow \text{Fin. coCartesian fibration}$ such that

$$\left\{ \begin{array}{l} \text{Ob}(\text{Corr}(C, E)^{\otimes}) \simeq \text{Corr}(C, E) \\ \text{s.m. structure so that } C^{\text{op}} \rightarrow \text{Corr}(C, E) \text{ is symm. monoidal} \\ \text{(so that in } \text{Corr}(C, E)) \\ X \otimes Y = X \times Y \end{array} \right.$$

$$\begin{array}{ccc} X & \xrightarrow{f} & X \\ \downarrow g & & \downarrow h \\ Y & & Y \end{array}$$

$$\begin{array}{ccc} X & \xrightarrow{f} & X \\ & & \searrow \\ & & Y \end{array}$$

DEF $\text{Corr}(C, E)^{\otimes} := \text{Corr}(C^{x'}, E^x)$

where $C^{x'} = ((C^{\text{op}})^{\cup})^{\text{op}}$, $C^{x'} \rightarrow \text{Fin.}^{\text{op}}$ (NB Not the s.m. category for (C, x))

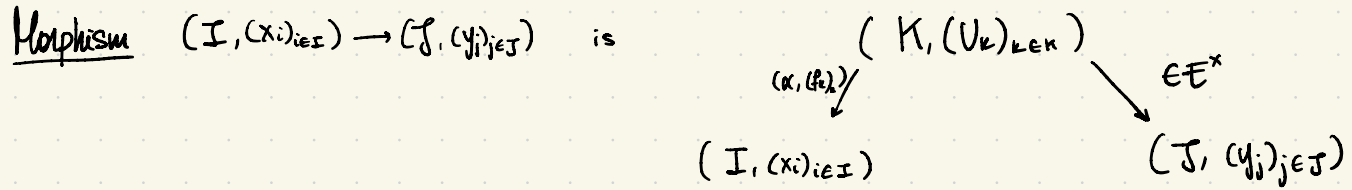
- $\text{Ob}(C^{x'}) = \text{Ob}((C^{\text{op}})^{\cup}) = \{ (c_i)_{i \in I} \mid I \in \text{Finset}, c_i \in C \}$

- Morphisms $(I, (x_i)_{i \in I}) \rightarrow (J, (y_j)_{j \in J})$ is $\left\{ \begin{array}{l} \alpha: \langle J \rangle \rightarrow \langle I \rangle \\ \forall i \in I \quad f_i: x_i \rightarrow \prod_{j \in \alpha^{-1}(i)} y_j \end{array} \right.$

and $E^x \subseteq \text{Morph}(C^{x'})$ lying over the identity

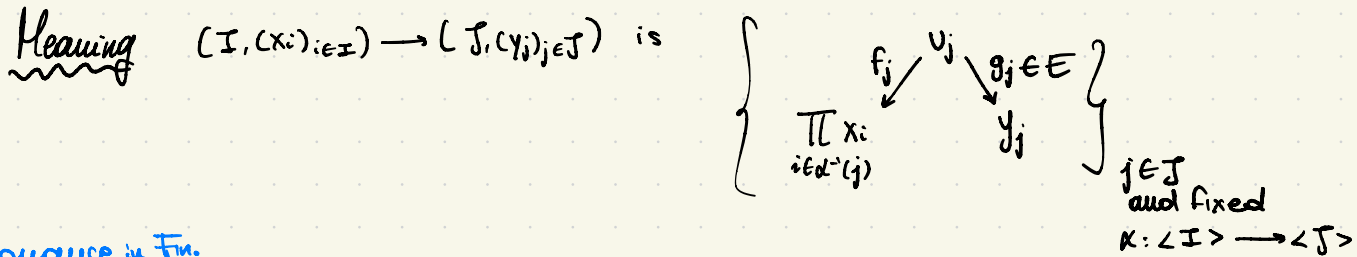
and s.t. $\forall i \in I \quad g_i: x_i \rightarrow y_i \in E$

So explicitly: $\text{Ob}(\text{Corr}(\mathcal{C}, \mathcal{E})^{\otimes}) = \{ (I, (x_i)_{i \in I}) \mid x_i \in \text{ob}(\mathcal{C}) \}$



So explicitly it is

$$\left\{ \begin{array}{l}
 \alpha: \langle I \rangle \rightarrow \langle K \rangle \\
 \forall k \in K \quad f_k: U_k \rightarrow \prod_{i \in \alpha^{-1}(k)} x_i \\
 \text{id}: \langle K \rangle = \langle J \rangle \\
 \forall k \in K \quad g_k: U_k \rightarrow y_k \in \mathcal{E}
 \end{array} \right.$$

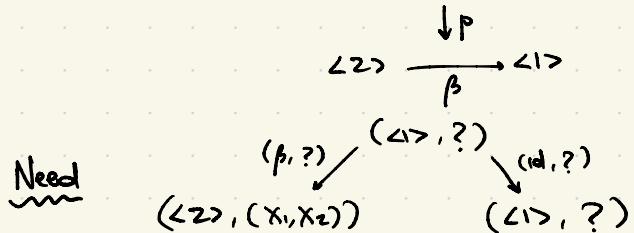


Rmk Contravariance in Fun.
to cancel
CONTRAVARIANCE GET down ✓

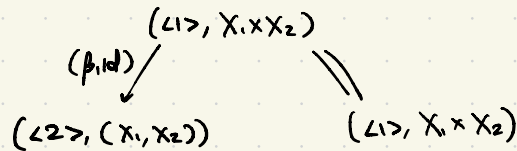
PROP [L2.2a] $\text{Corr}(C, E)^{\otimes} \xrightarrow{p} \text{Fin.}$ is a COCAUSTESIAN FIBRATION, and $\text{Corr}(C, E)^{\otimes}_{\langle 1 \rangle} \simeq \text{Corr}(C, E)^{\otimes \times 1}$

Rmk The s.m. structure on $\text{Corr}(C, E) \simeq \text{Corr}(C, E)^{\otimes}_{\langle 1 \rangle}$ is the coCat left of $\beta: \langle 2 \rangle \rightarrow \langle 1 \rangle$

Consider $(\langle 2 \rangle, (x_1, x_2)) \dashrightarrow ?$



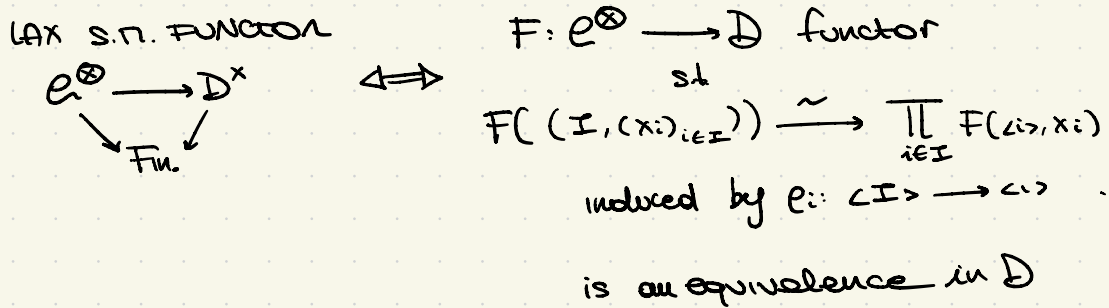
and this is



2.6. A 3-FUNCTOR FORMALISM

We said that a 3-functor formalism is LAX S.T. FUNCTOR $\text{Corr}(\mathcal{C}, \mathcal{E})^{\otimes} \longrightarrow \text{Cat}_{\infty}^{\times}$

Thm [2.6.1.7, HA] $\mathcal{E}^{\otimes} \rightarrow \text{Fun. s.m. cocat}$
 $\mathcal{D}^{\times} \rightarrow \text{Fun. CARTESIAN s.m. cocat}$



Rmk Easy to see \Leftarrow , bc $F(\prod_i X_i) \xleftarrow{F(\beta)} F((I, (x_i)_{i \in I})) \xrightarrow{\sim} \prod_{i \in I} F(\langle i \rangle, x_i)$

β cocat lift of $\beta_I: \langle I \rangle \longrightarrow \langle I \rangle$

DEF | (C, E) geometric context. Then a 3-FUNCTOR FORMALISM is a functor

$$D: \text{Corr}(C, E)^{\otimes} \longrightarrow \text{Cat}_{\infty} \quad \text{s.t.} \quad D(I, (X_i)_{i \in I}) \xrightarrow{\sim} \prod_{i \in I} D(\langle i \rangle, X_i)$$

is an equivalence.

In practice, objectwise D is determined by the obj's

$$\Rightarrow \text{we write } D(X) := D(\langle 1 \rangle, X)$$

$$\text{Corr}(C, E)_{\langle 1 \rangle}^{\otimes} \simeq \text{Corr}(C, E)$$

$$(\langle 1 \rangle, Y) \longleftrightarrow Y$$

RESULT: D a 3 functor formalism \rightsquigarrow we recover the 3 functors $\otimes, f^*, f!$
 and FUNCTIONALITY + COMPATIBILITY encoded by functoriality of D
 (• PROJ FORMULA
 • BASE CHANGE)

Indeed

• PULLBACK

$f: X \rightarrow Y$ morph in \mathcal{C}

$$\Rightarrow \begin{array}{ccc} & (\langle 1 \rangle, X) & \\ \text{(id, f)} \swarrow & & \parallel \\ (\langle 1 \rangle, Y) & & (\langle 1 \rangle, X) \end{array}$$

$$\text{morph } (\langle 1 \rangle, Y) \xrightarrow{f} (\langle 1 \rangle, X)$$

$$\rightsquigarrow f^* = D(f): D(Y) \longrightarrow D(X)$$

• TENSOR PRODUCT

Want $(D(X), \otimes)$ s.m. category.

Let's see how to define $-\otimes-$: $D(X) \times D(X) \longrightarrow D(X)$

EXTERNAL
TENSOR
PRODUCT

$$1) \boxtimes: D(X) \times D(X) \xleftarrow{\sim} D(\langle 2 \rangle, (X, X)) \xrightarrow{D(\beta)} D(X \times X)$$

$$2) \text{diag}: X \rightarrow X \times X \Rightarrow \text{diag}^*: D(X \times X) \rightarrow D(X)$$

$$\Rightarrow -\otimes-: D(X) \times D(X) \xleftarrow{\sim} D(\langle 2 \rangle, (X, X)) \xrightarrow{D(\beta)} D(X \times X) \xrightarrow{\text{diag}^*} D(X)$$

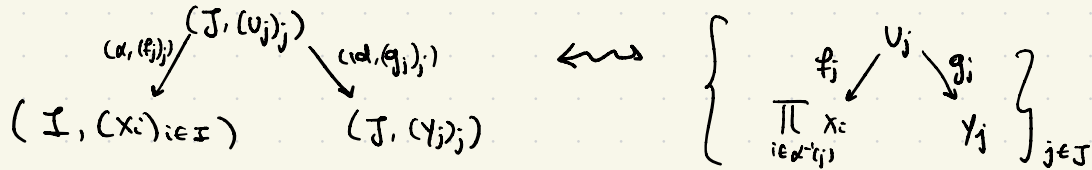
• EXCEPTIONAL PUSHFORWARD

$$f: X \rightarrow Y \text{ in } \mathcal{E} \Rightarrow \begin{array}{ccc} & (\langle 1 \rangle, X) & \\ \parallel & & \searrow \text{(id, f)} \\ (\langle 1 \rangle, X) & & (\langle 1 \rangle, Y) \end{array}$$

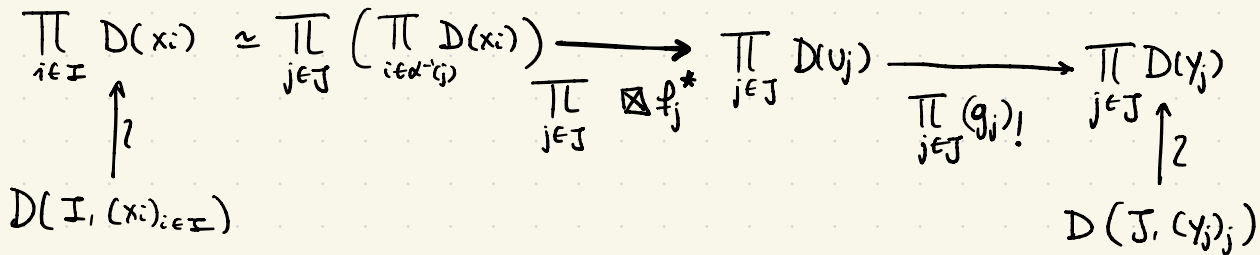
$$\rightsquigarrow f_! = D(f): D(X) \longrightarrow D(Y)$$

Compatibility laws \iff functoriality of D

• ACTION of D on a generic morphism:



and



Some compatibility relation

PROJECTION FORMULA: $f! A \otimes B \simeq f! (A \otimes f^* B)$ functionally

$f: X \rightarrow Y$ in \mathcal{E}

i.e. $\underbrace{f!(-) \otimes -}_{(1)} \simeq \underbrace{f!(- \otimes f^*(-))}_{(2)}$ as functors $\mathcal{D}(X) \times \mathcal{D}(Y) \rightarrow \mathcal{D}(Y)$

$$f! (f! \beta) = \beta \cdot \mathcal{D} \left(\begin{array}{ccc} & (\langle 1 \rangle, Y \times Y) & \\ \beta \cdot \text{id} \swarrow & \parallel & \text{diag} \swarrow \\ (\langle 2 \rangle, (Y, Y)) & (\langle 1 \rangle, Y \times Y) & (\langle 1 \rangle, Y) \end{array} \right)$$

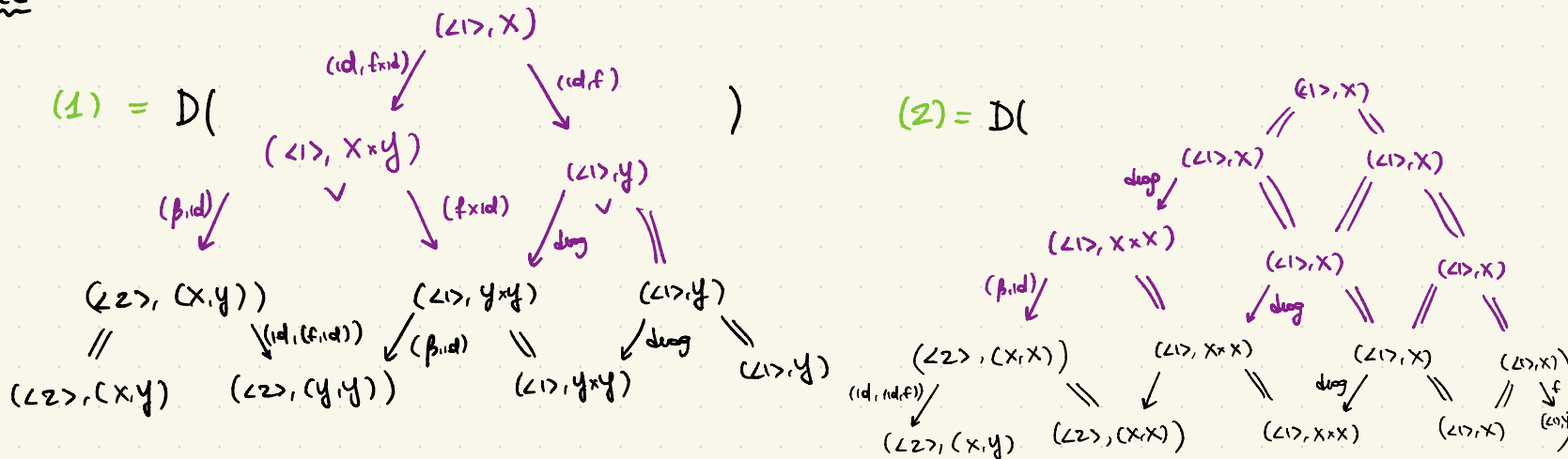
$$(2) = \mathcal{D} \left(\begin{array}{ccccccc} & (\langle 2 \rangle, (X, X)) & & (\langle 1 \rangle, X \times X) & & (\langle 1 \rangle, X) & & (\langle 1 \rangle, X) \\ \text{id} \cdot \text{id} \cdot f \swarrow & \parallel & \cdot & \swarrow & \parallel & \cdot & \text{diag} \swarrow & \parallel & \cdot & \parallel & \downarrow f \\ (\langle 2 \rangle, (X, Y)) & (\langle 2 \rangle, (X, X)) & & (\langle 2 \rangle, (X, X)) & & (\langle 1 \rangle, X \times X) & & (\langle 1 \rangle, X \times X) & & (\langle 1 \rangle, X) & & (\langle 1 \rangle, X) \end{array} \right)$$

to read in this direction \rightarrow

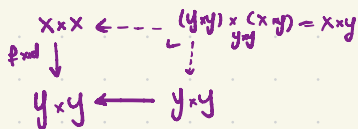
Now if we look at the compositions of the two morphisms in $\text{Com}(C, E)^\otimes$

They are the SAME!

See



where first pullback for example is



$$\Rightarrow D\left(\begin{array}{c} \text{(\langle 1 \rangle, X)} \\ \swarrow \text{(\beta, id \circ f)} \quad \searrow \text{(id, f)} \\ \text{(\langle 2 \rangle, X \times Y)} \quad \text{(\langle 1 \rangle, Y)} \end{array} \right) \quad \underline{\underline{OK}}$$

Similar way we get BASE CHANGE.

Final Remark: $f: X \rightarrow Y$ in \mathcal{E} \rightsquigarrow $D(Y) \curvearrowright D(X)$ via $D(Y) \times D(X) \xrightarrow{f^* \times \text{id}} D(X) \times D(X) \xrightarrow{\otimes} D(X)$

and WANT (in a 3 fact formalism) that $\#!: D(X) \rightarrow D(Y)$ is $D(Y)$ -lin.

This is TWE bc D is LAX S.N. functor

Motivation: • D lax s.m. \rightarrow D preserves algs & mods

$$\begin{array}{ccc} \cdot (C^{\text{op}})^U & \longrightarrow & \text{Corr}(C, E)^{\otimes} \\ \downarrow f \text{ in } C & & \downarrow f \text{ in } C \end{array} \quad \text{s.m. functor}$$

X is a module over y in C^{op}
(as comm algs in (C^{op}, U))

$$\begin{array}{ccc} X \cup X (= X \times X) & \xrightarrow{\Delta} & X \\ \uparrow f \cup f & \searrow \eta & \uparrow f \\ Y \cup Y & \xrightarrow{\Delta} & Y \end{array}$$

and $g: X \rightarrow X'$ map / y in E

$$\begin{array}{ccc} X & \xrightarrow{g} & X' \\ f \searrow & & \swarrow f' \end{array}$$

\Downarrow
induces
 $X \rightarrow X'$ map of y -mod in $\text{Corr}(C, E)$

(LAX S.N.)
 $\Rightarrow D(g): g!: D(X) \rightarrow D(X')$ \rightsquigarrow Now take $X' = Y$.
map of $D(Y)$ -mod